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# **HPC-FF – Overview and Experience**

**eGOTiT Seminar, 15<sup>th</sup> Dec**

**Nitya Hariharan**  
**High Level Support Team**  
**Max-Planck-Institut für Plasmaphysik, München, Germany**



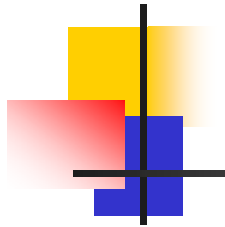
**With contributions from:**  
**IPP, Germany; University of Alicante, Spain**



# Outline

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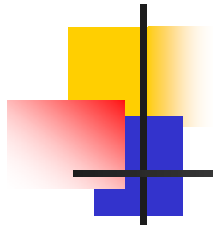
- **HPC-FF**
- **Intel Nehalem Architecture**
- **Infiniband network**
- **BEUPACK Benchmark**
- **Shared Memory Segments**
- **CPU affinity**



# HPC-FF

- **Start of production – Aug 2009**
- **1080 nodes @ 8 cores,  
peak performance 101 TFlop/s**
- **Quad core Intel Xeon 5570 (Nehalem)  
processors**
- **24 Gb per node, Lustre filesystem**
- **Infiniband network with Mellanox switch**
- **ParTec MPI, Intel MPI**





# Intel Nehalem

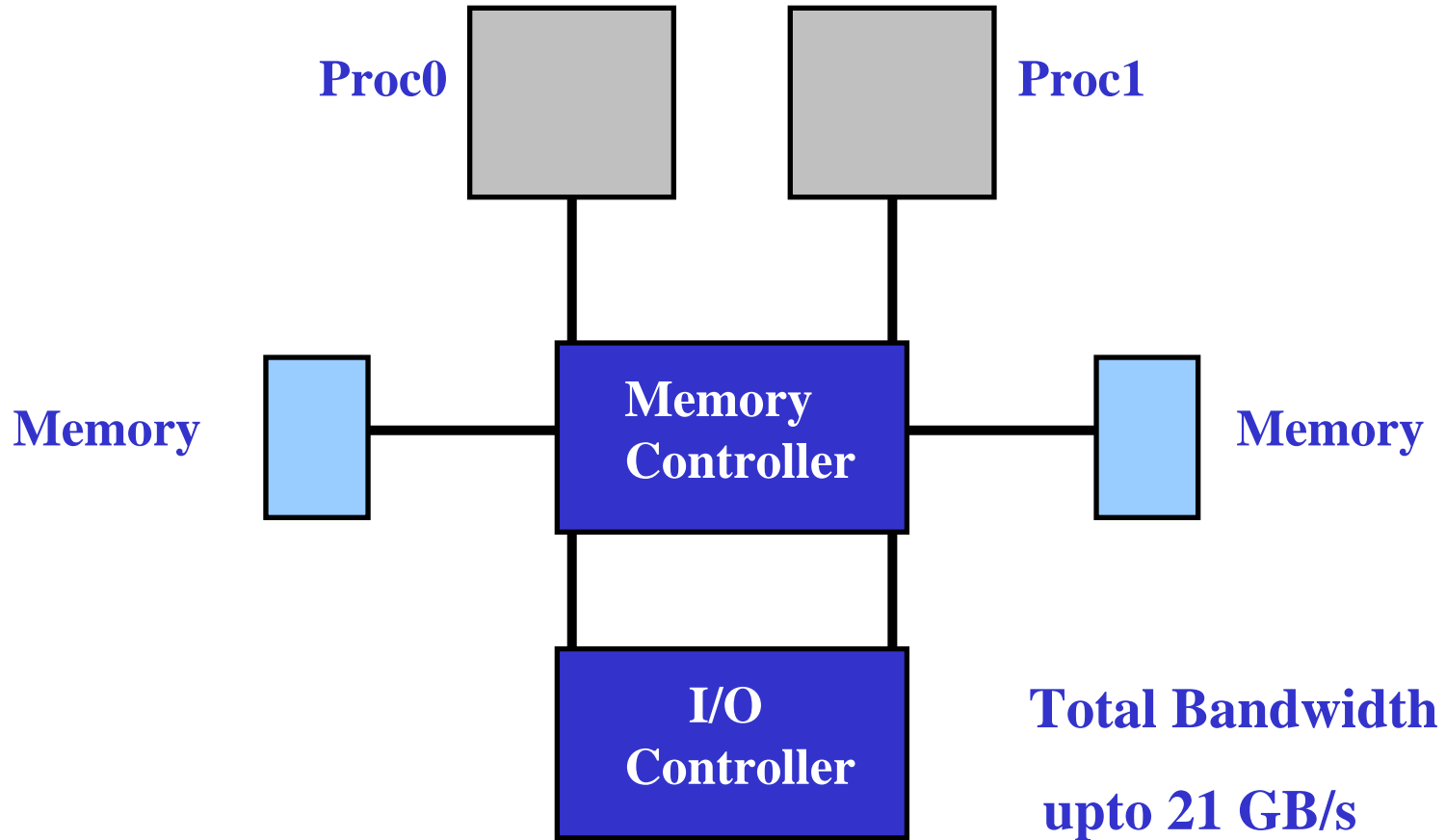
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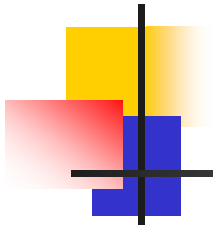
- Supports SSE4.2 instructions
- QPI – Quick Path Interconnect
- Theoretical **raw** bandwidth – 25.6 GB/s
- Dynamic resource scaling
- NUMA architecture - low latency for local memory

Source - Intel

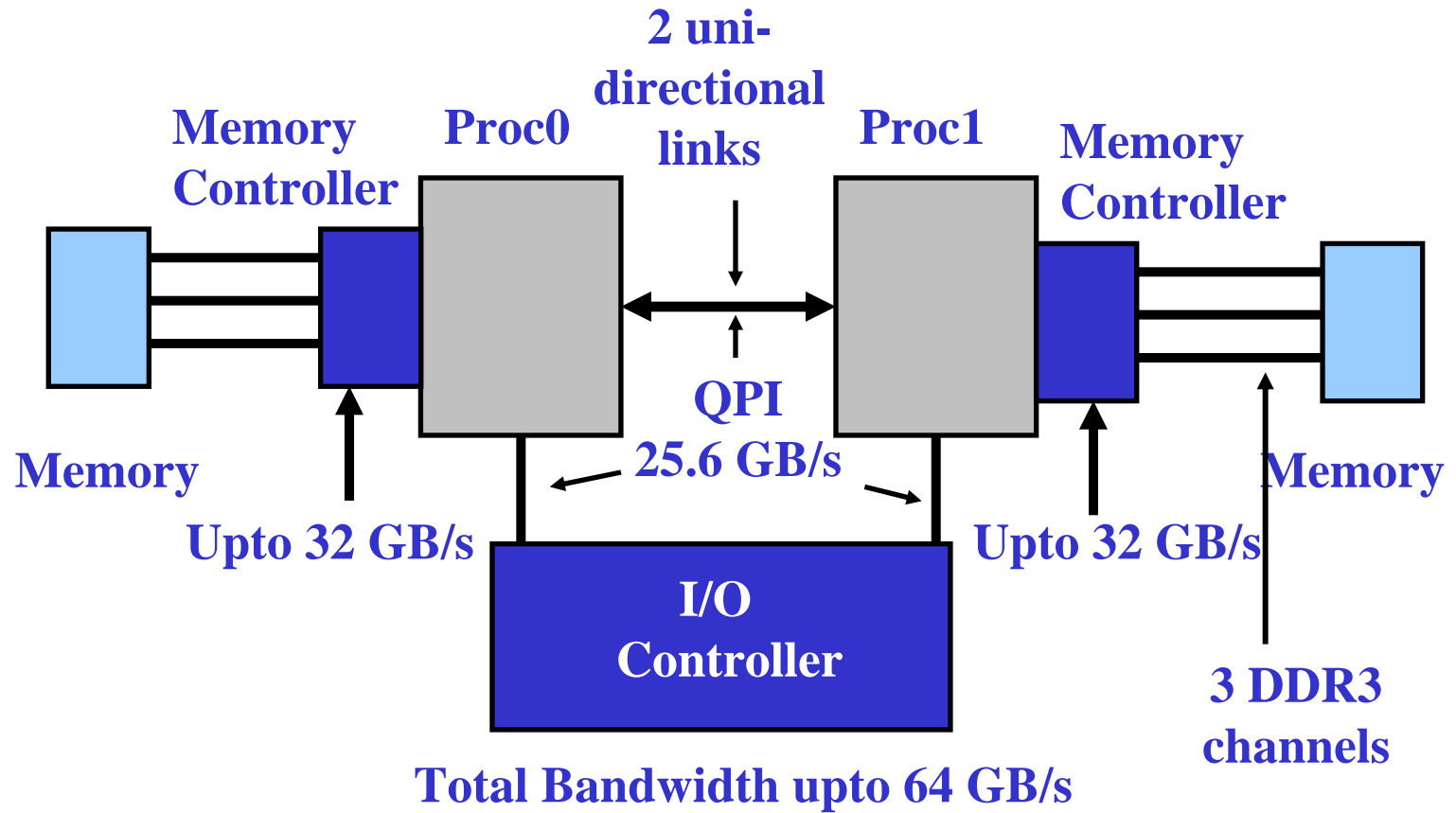


# Before QPI



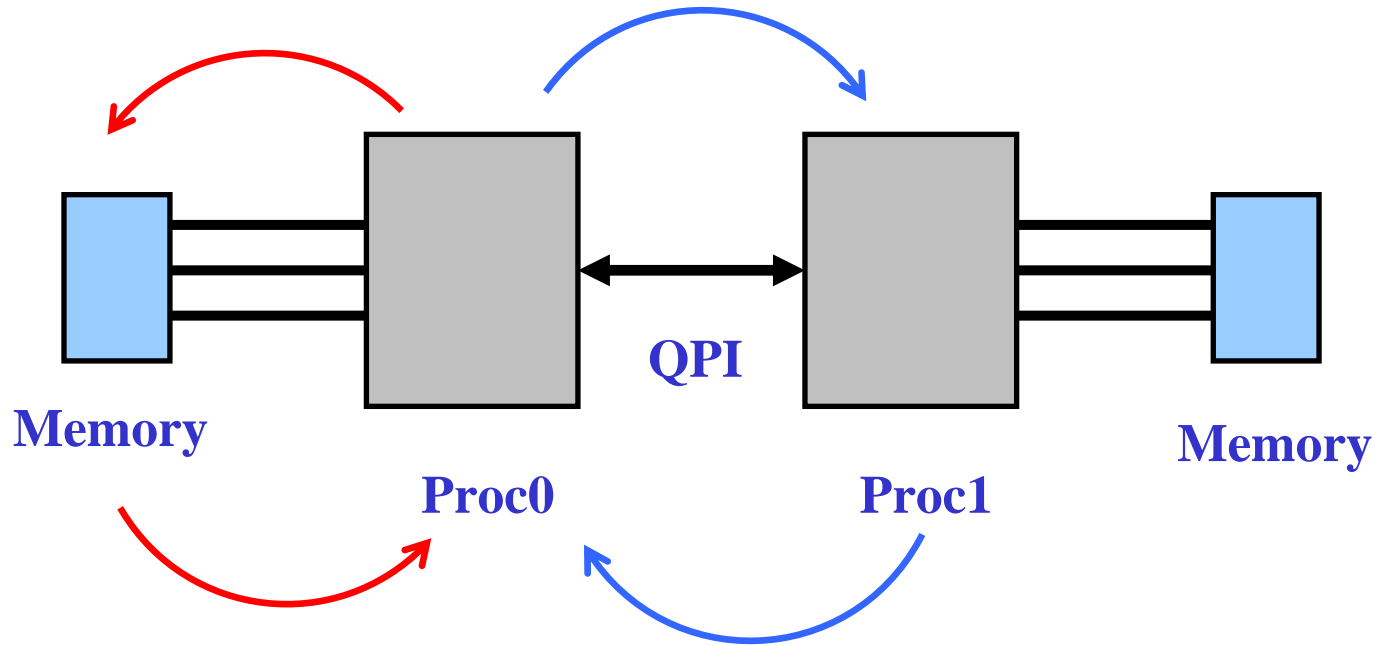


# QPI

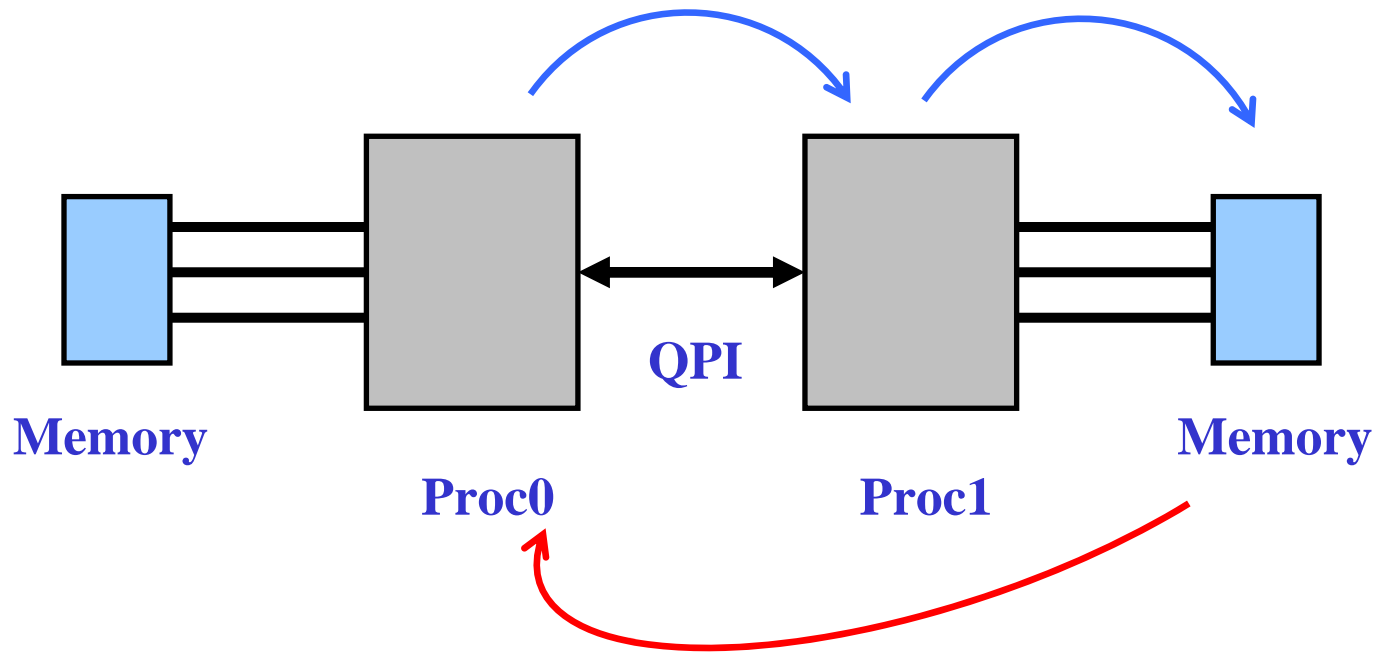




# Local Memory access



# Remote Memory access



Source - Intel

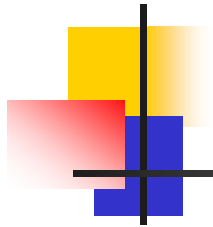




# Infiniband(IB) Network

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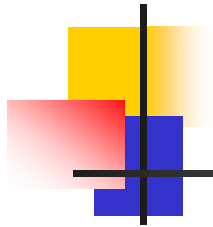
- **Fabric/network topology for interconnecting compute and I/O nodes**
- **Nodes connected to the network through Host Channel Adaptors (HCAs)**
- **Queue based model, Send and Receive Queue (Queue Pair – QP)**
- **RDMA (Remote Direct Memory Access) also supported**



# RDMA

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- **Direct access to memory location of remote processes**
- **Uses zero-copy mechanism – origin and destination buffers are registered prior to use**
- **Virtual address and memory access key of remote process reqd**
- **Lower end-to-end latencies**
- **RDMA Read and Write**



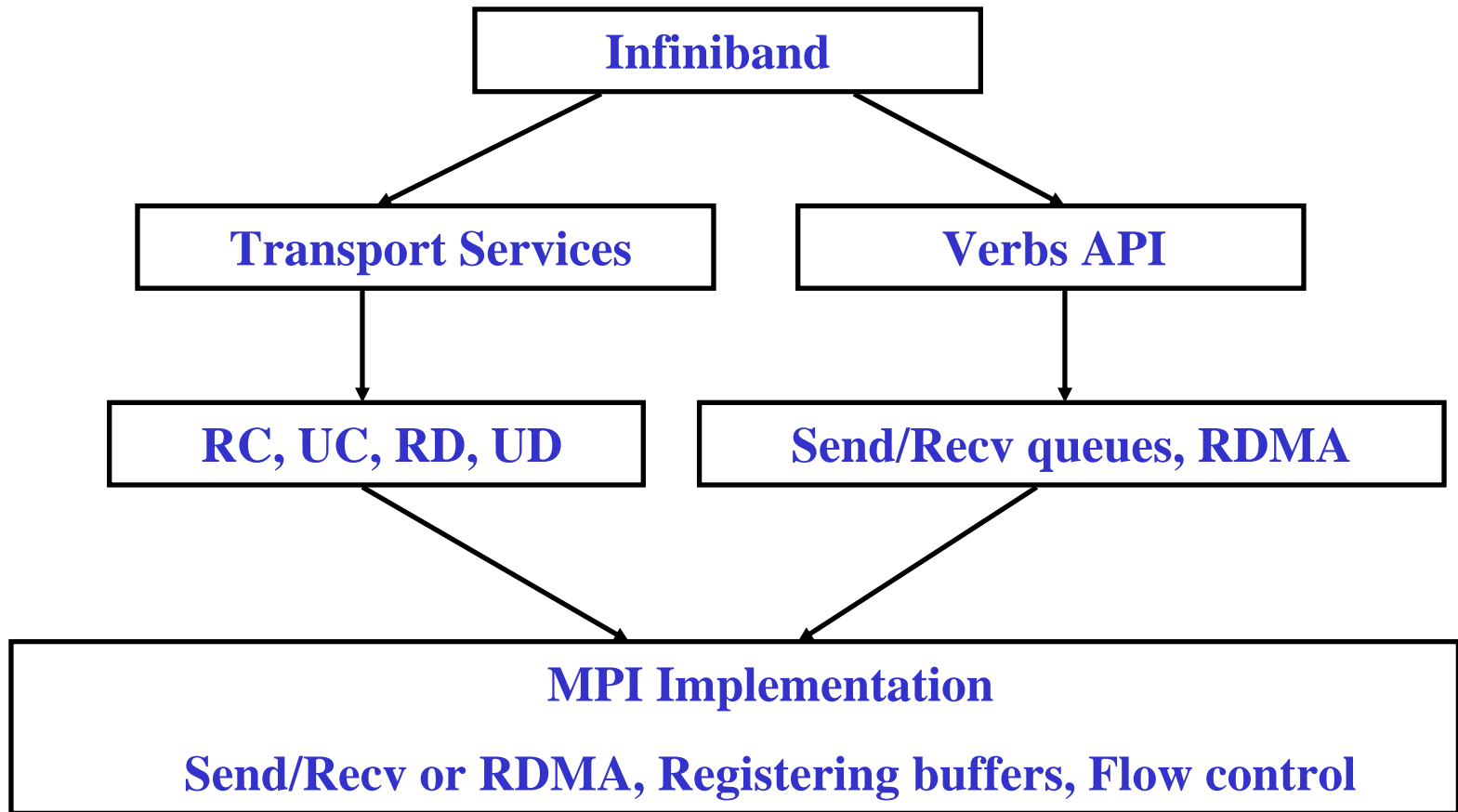
# Infiniband

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- **Different types of transport service**
  - ✓ **Reliable connection**
  - ✓ **Unreliable connection**
  - ✓ **Reliable datagram**
  - ✓ **Unreliable datagram**
- **On HPC-FF, RDMA writes with reliable connections as of now**
- **Beta implementation of Unreliable datagram**



# IB and MPI





# BEUPACK Benchmark

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**Fusion codes used as benchmark for IFERC procurement**

- **ORB5**
- **GENE**
- **GEMR**
- **JOREK**
- **MDCASK**
- **GYSELA**



# ORB5 benchmark

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- **Strong scaling –  $2048 \times 10^6$  ions**
- **Weak scaling –  $256 \times 10^6$  ions (256 cores) and then quadrupled**
- **Memory issues on 4096 cores, used PSP\_ONDEMAND**

**Test cases by A. Bottino**



# PSP\_ONDEMAND

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- ParTec MPI requires 0.55 MB per connection
- 16 buffers for Send and Receive
- ~2.2 GB per core for 4096 cores -> 500 MB left for application
- PSP\_ONDEMAND=1, dynamic memory allocation

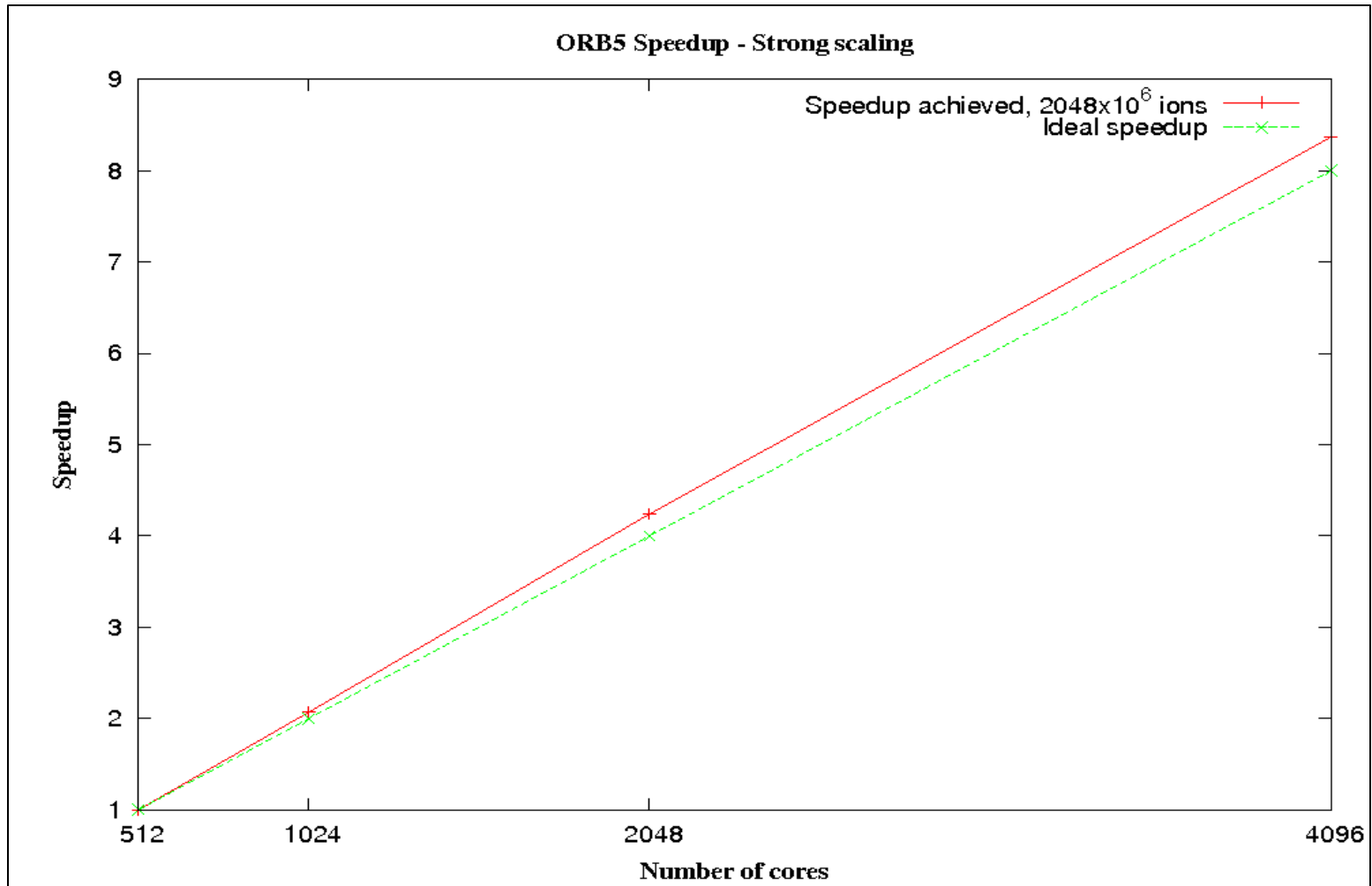
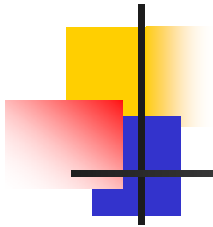


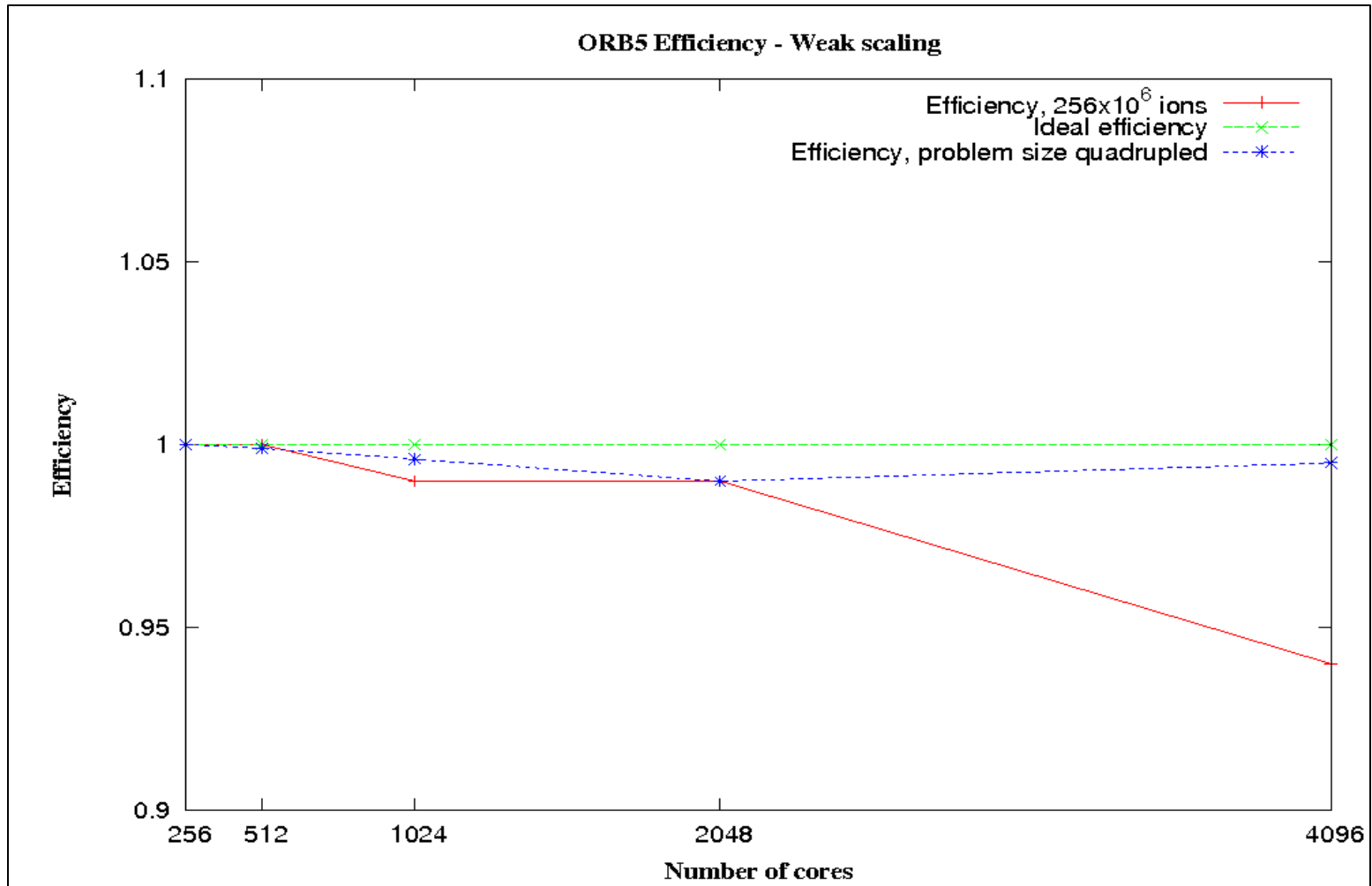
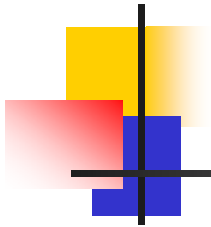
# PSP\_ONDEMAND

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- Works well as long as no all-to-all communication or large core numbers
- `PSP_OPENIB_SENDQ_SIZE` and `PSP_OPENIB_RECVQ_SIZE`
- Can degrade MPI's throughput and messaging rate, should be  $\geq 3$  to prevent deadlocks.









# GENE benchmark

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- **Strong scaling – 1.752 GB for 512 cores**
- **Weak scaling – 1.752 GB (constant)**
- **Usage of MPI\_ANY\_SOURCE**

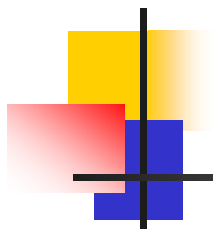
**Test cases by T. Dannert**



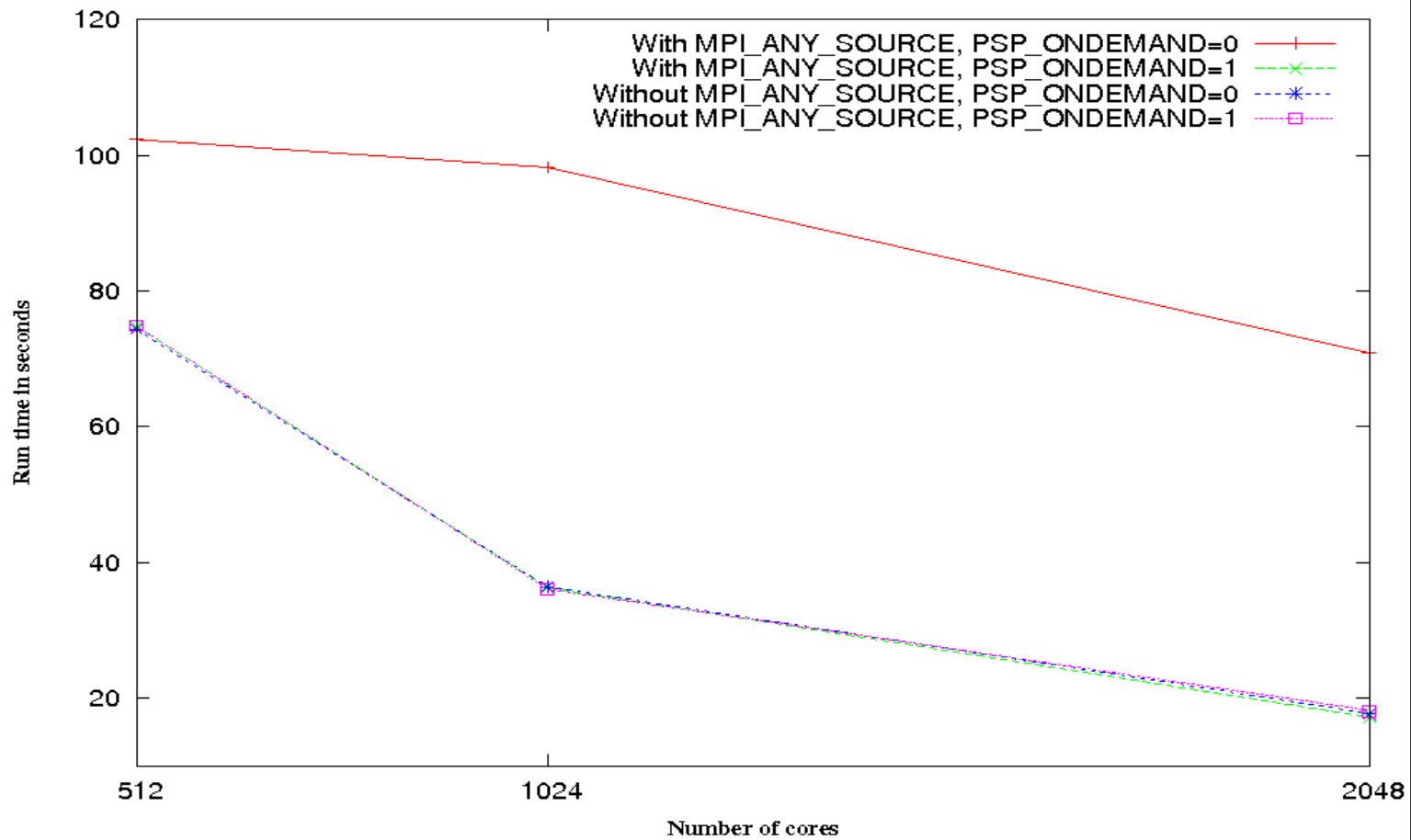
# MPI\_ANY\_SOURCE

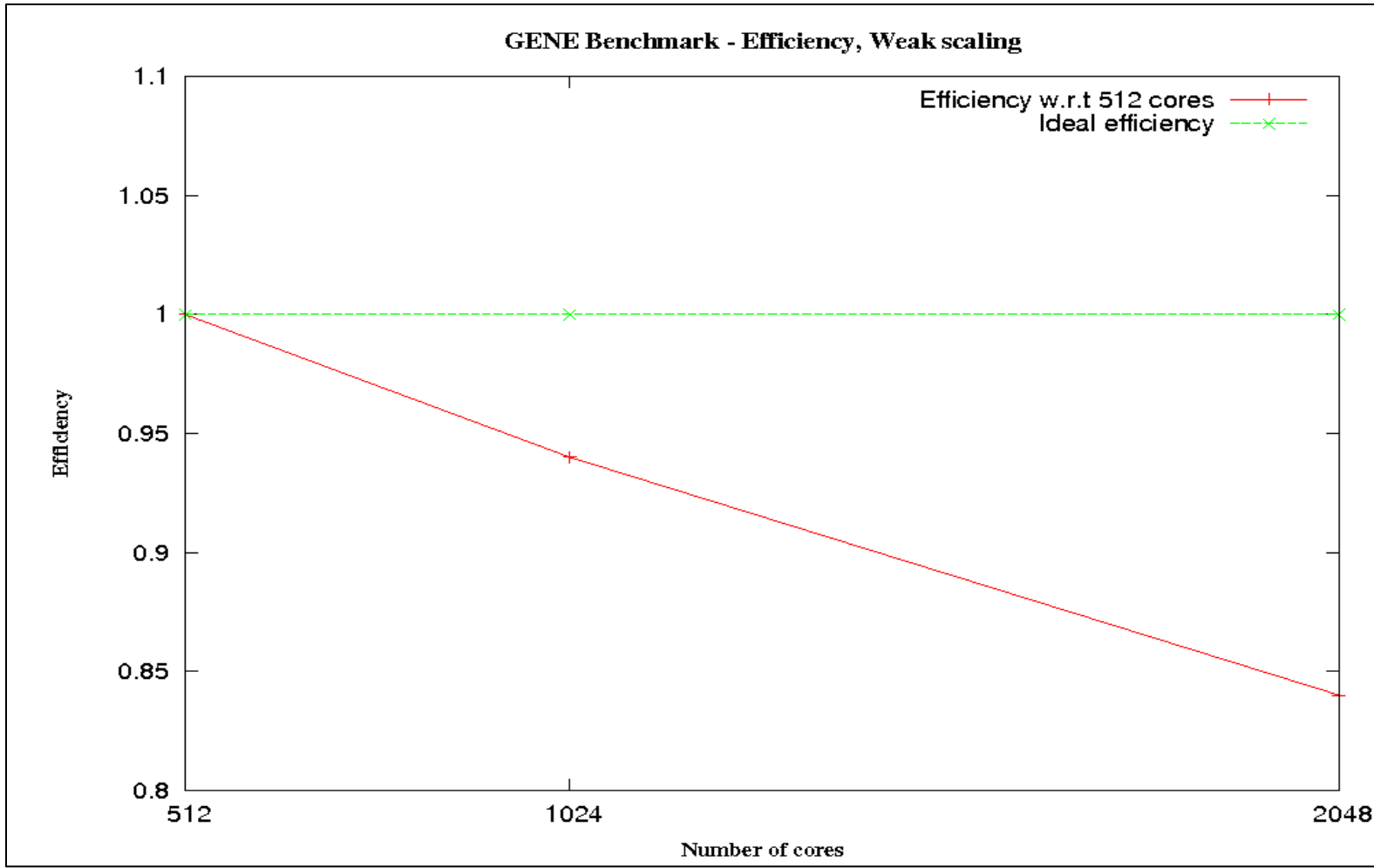
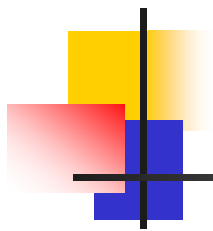
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- **Communication lib has to poll a long list of receive queues**
- **Avoid it!**
- **When PSP\_ONDEMAND=1, only one File Descriptor (FD) has all unestablished connections**
- **Few IB buffers from established connections and one FD**



GENE Benchmark - Strong scaling, with/without MPI\_ANY\_SOURCE





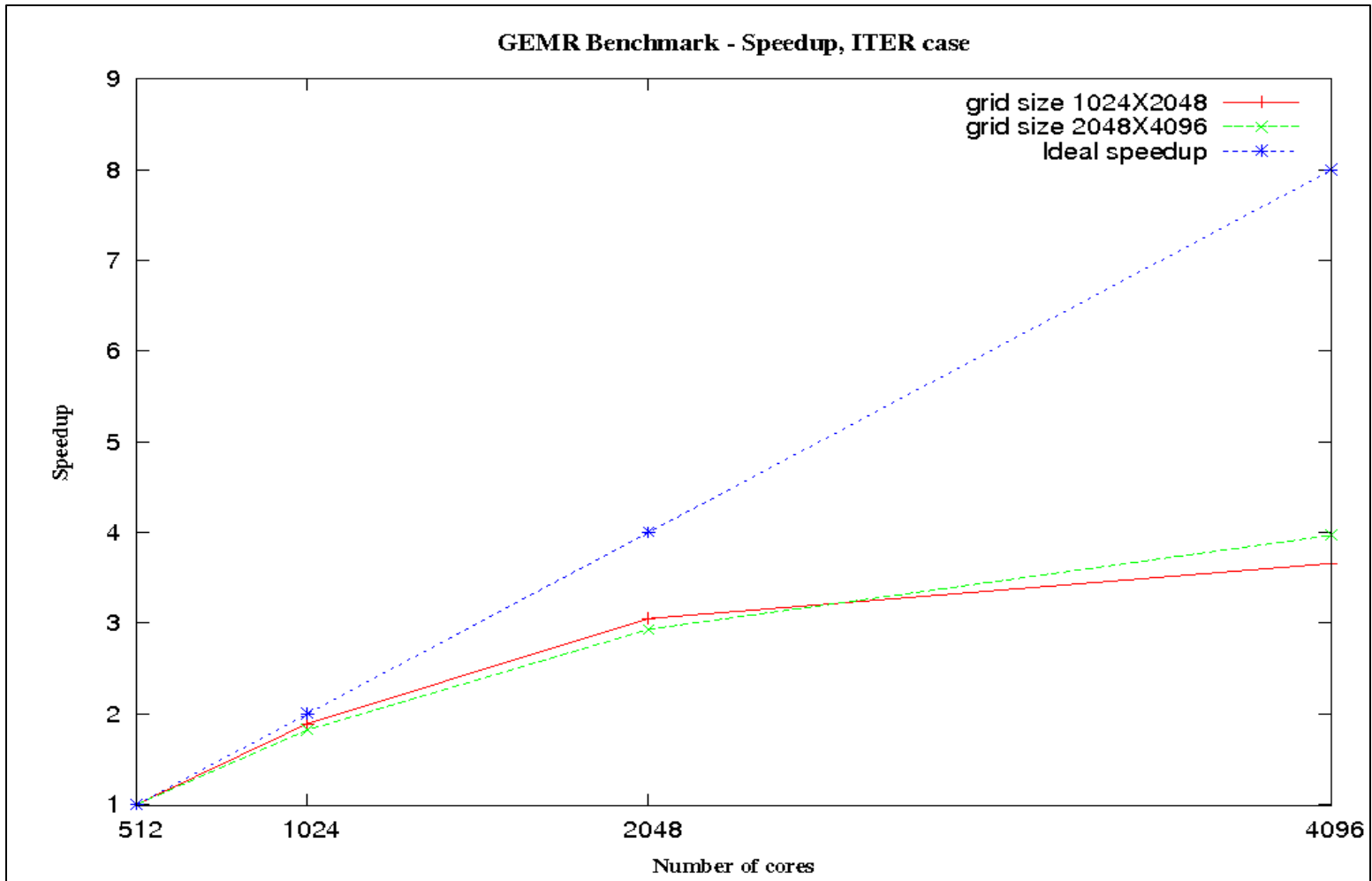
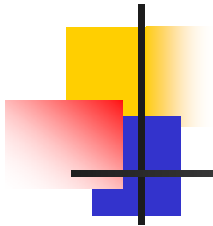


# GEMR benchmark

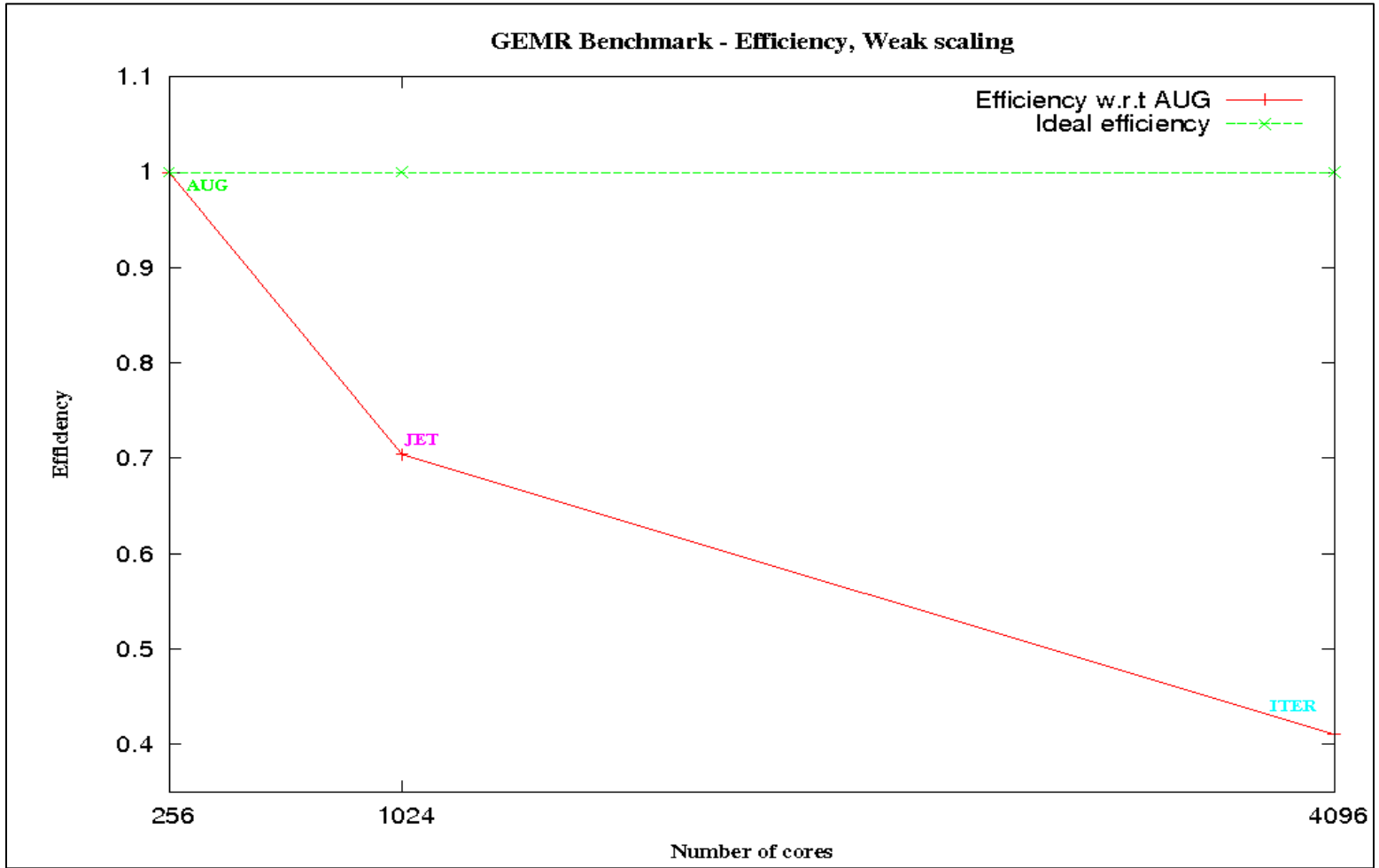
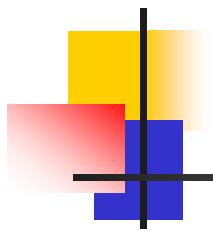
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- **Strong scaling – ITER cases, grid size 1024x2048, 2048x4096**
- **Weak scaling – AUG, JET and ITER cases**
  - ✓ **AUG – 256x512**
  - ✓ **JET – 512x1024**
  - ✓ **ITER – 1024x2048**
- **No system related issues**

**Test cases by B.D. Scott**









# JOREK benchmark

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- **Only hybrid code as part of the benchmark**
- **Uses PastiX library – parallel solver for large sparse systems**
- **MPI\_THREAD\_MULTIPLE (MTM)**
- **MPI\_THREAD\_FUNNELED**
  - **errors with MPI\_Start/MPI\_Startall**
- **Test runs on IBM AIX, Sun Linux cluster AIMS (Intel MPI)**
- **Intel MPI installed on HPC-FF - /usr/local/impi/3.2.2.006/bin64/**
- **Working version of ParTec MPI with MTM support available**

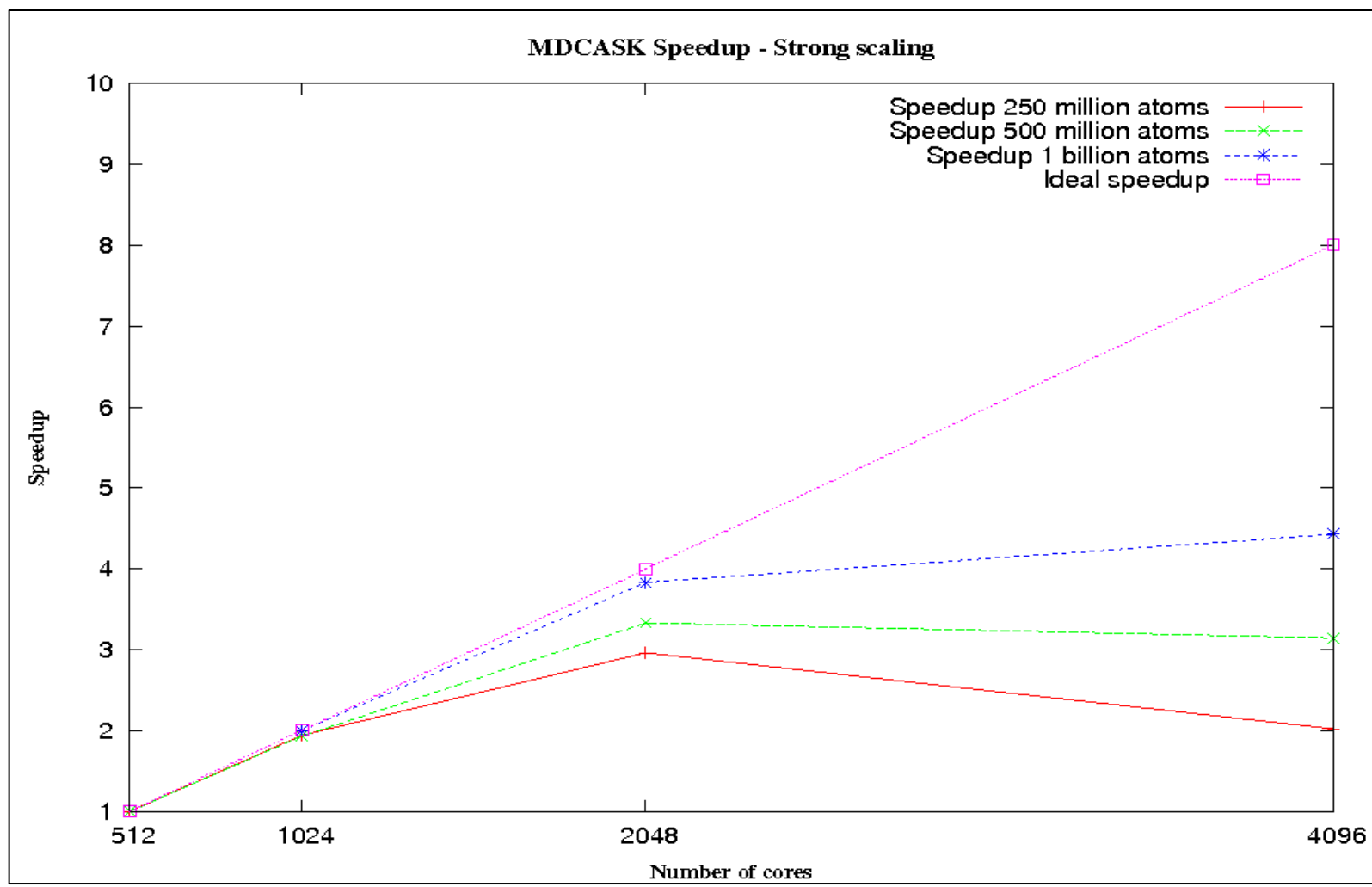
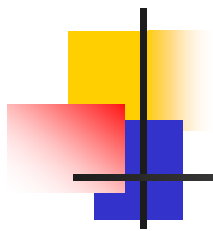


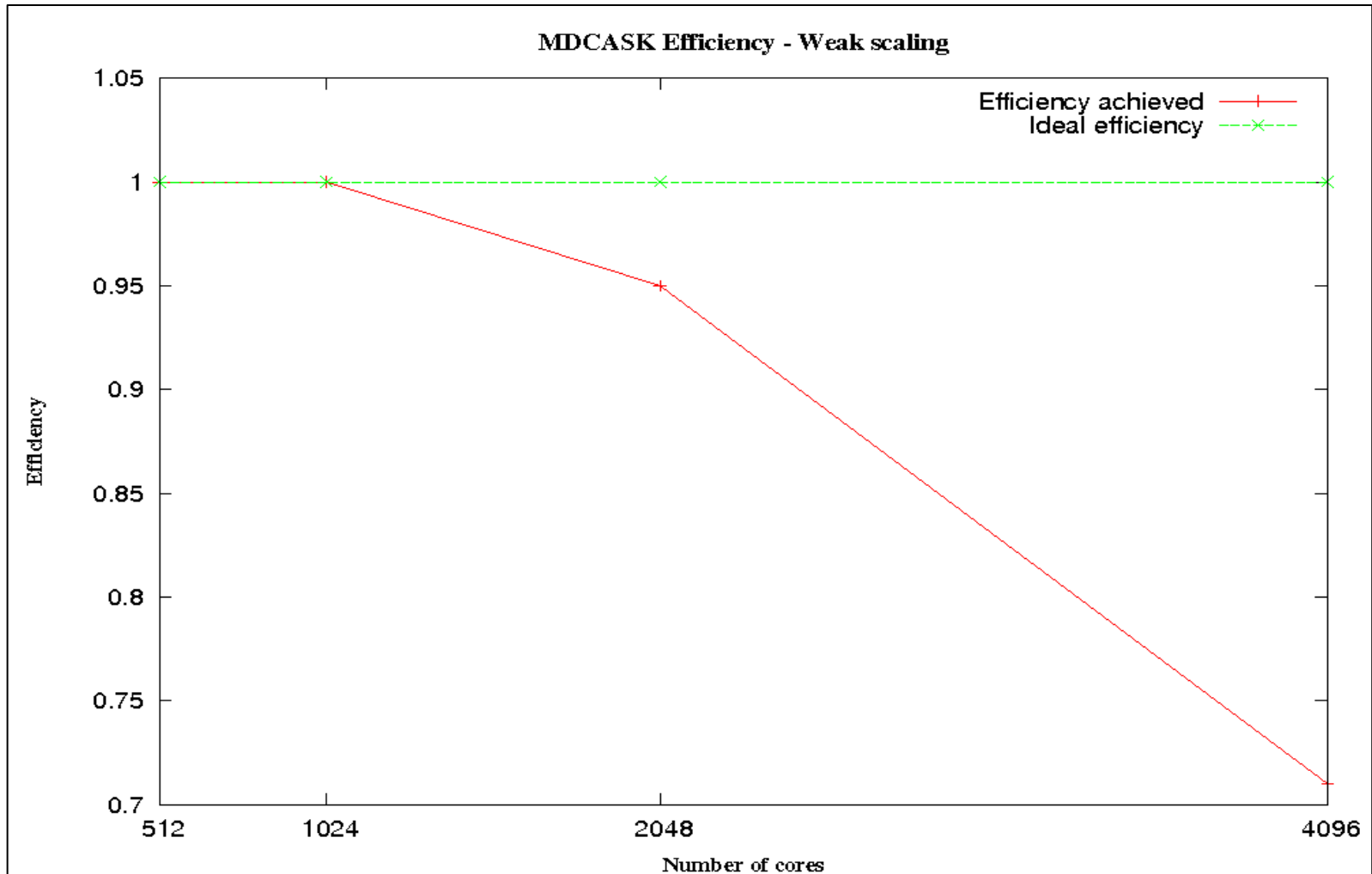
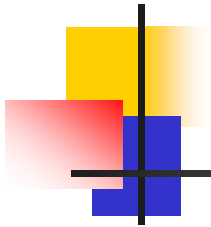
# MDCASK benchmark

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- **Strong scaling – upto 1 billion atoms, too large for 512 cores (memory issues)**
- **Weak scaling – upto 2 billion atoms**
- **PSP\_ONDEMAND – Intel MPI**
- **Intel does dynamic buffer allocation by default  
I\_MPI\_USE\_DYNAMIC\_CONNECTIONS**

**Test cases by M.J. Caturia**



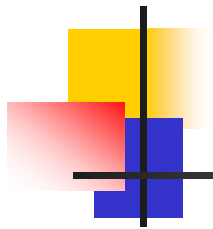




# ParTec and Intel MPI

Memory consumption with static buffer allocation

Number of cores	ParTec MPI	Intel MPI
128	68 MB	68 MB
256	138 MB	139 MB
512	278 MB	281 MB
1024	558 MB	565 MB
2048	1118 MB	Program crash

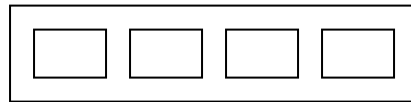


# PSI\_TPP

- NUMA design
- Memory separated into two nodes of 12 GB each

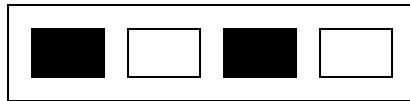


**Processor 0**

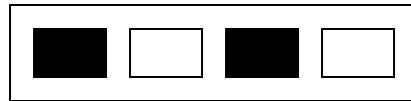


**Processor 1**

**Default  
(PSI\_TPP = 1)**



**Processor 0**



**Processor 1**

**PSI\_TPP = 2**



# PSI\_TPP

<b>ppn</b>	<b>PSI_TPP</b>	<b>Number of cores per compute node</b>	<b>Memory per core</b>	<b>Explanation</b>
<b>8</b>	<b>1 (default)</b>	<b>8</b>	<b>3</b>	<b>1 core allocated per task.</b>
<b>8</b>	<b>2</b>	<b>4</b>	<b>6</b>	<b>2 cores allocated per task, 2*3GB memory available per core. Doubled bandwidth.</b>
<b>8</b>	<b>4</b>	<b>2</b>	<b>12</b>	<b>4 cores allocated per task.</b>
<b>4</b>	<b>1 (default)</b>	<b>4</b>	<b>3</b>	<b>1 core allocated per task.</b>





# PSI\_TPP

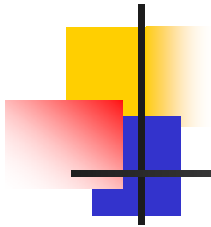
<b>Benchmark (512 cores)</b>	<b>Run time in seconds (PSI_TPP=1)</b>	<b>Run time in seconds (PSI_TPP=2)</b>	<b>%Gain</b>
<b>ORB5</b>	<b>309.34</b>	<b>280.12</b>	<b>9.45</b>
<b>MDCASK – 250 Million atoms</b>	<b>4.80 (per time step)</b>	<b>4.64 (per time step)</b>	<b>3.33</b>
<b>MDCASK – 500 Million atoms</b>	<b>9.26 (per time step)</b>	<b>8.82 (per time step)</b>	<b>4.75</b>
<b>GEMR – ITER case (problem size 4096x2048)</b>	<b>2730.28</b>	<b>1985.62</b>	<b>27.27</b>
<b>GEMR – ITER case (problem size 2048x1024)</b>	<b>682.37</b>	<b>428.7</b>	<b>37.17</b>
<b>GENE</b>	<b>74.87</b>	<b>49.99</b>	<b>33.23</b>



# Shared Memory Segments (SMS)

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- **SMS – area of memory accessible by more than one process on the same node**
- **Based on System V IPC (Inter Process Communication) framework**
- **FIPC (Fortran IPC) module by Ian Bush**
- **Fortran and C interoperability**
- **Hybrid codes – MPI and OpenMP, effort in adding OpenMP to existing MPI codes**
- **Useful for machines like HPC-FF**
- **More portable, no thread-safe MPI reqd**



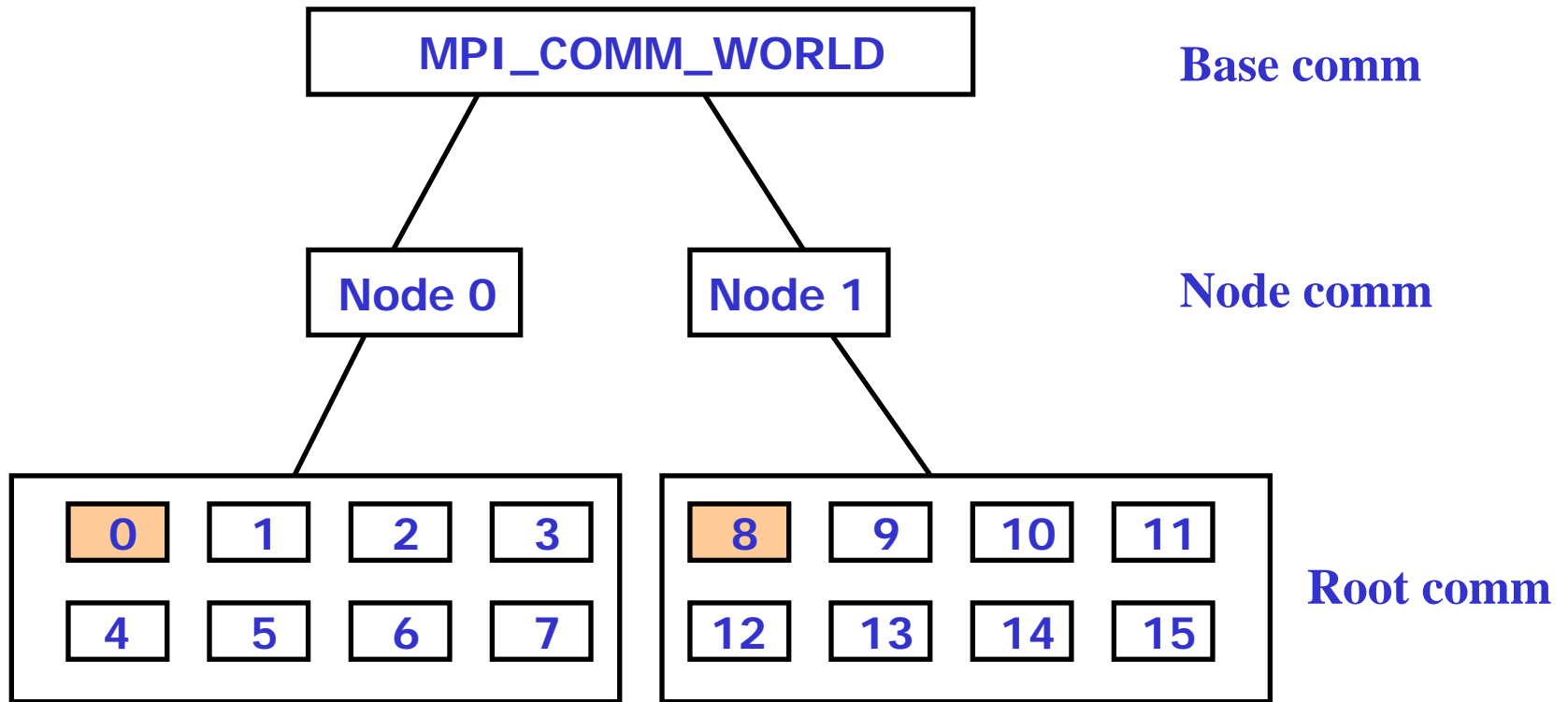
# FIPC

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- Create a context – *fipc\_init*
- Create a SMS in this context – *fipc\_seg\_create*
- Processes with an SMP node share the memory segment
- Need semaphores – *fipc\_critical\_start/end*
- Routines like *fipc\_allreduce*, *fipc\_ctxt\_rank*, *fipc\_ctxt\_size*



# FIPC Context





# Splines and SMS

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- Tested with Cubic Splines from Numerical Recipes
- Two functions *spline* and *splint*
- Input data  $x, f(x)$  and output of *spline* can be used as SMS.
- Will be used for EZspline library for ASCOT.
- Downside, SMS not deleted after abnormal termination. Need additional processing. On batch nodes, cleaned up by ParTec software, but can still be a problem with multiple jobs.



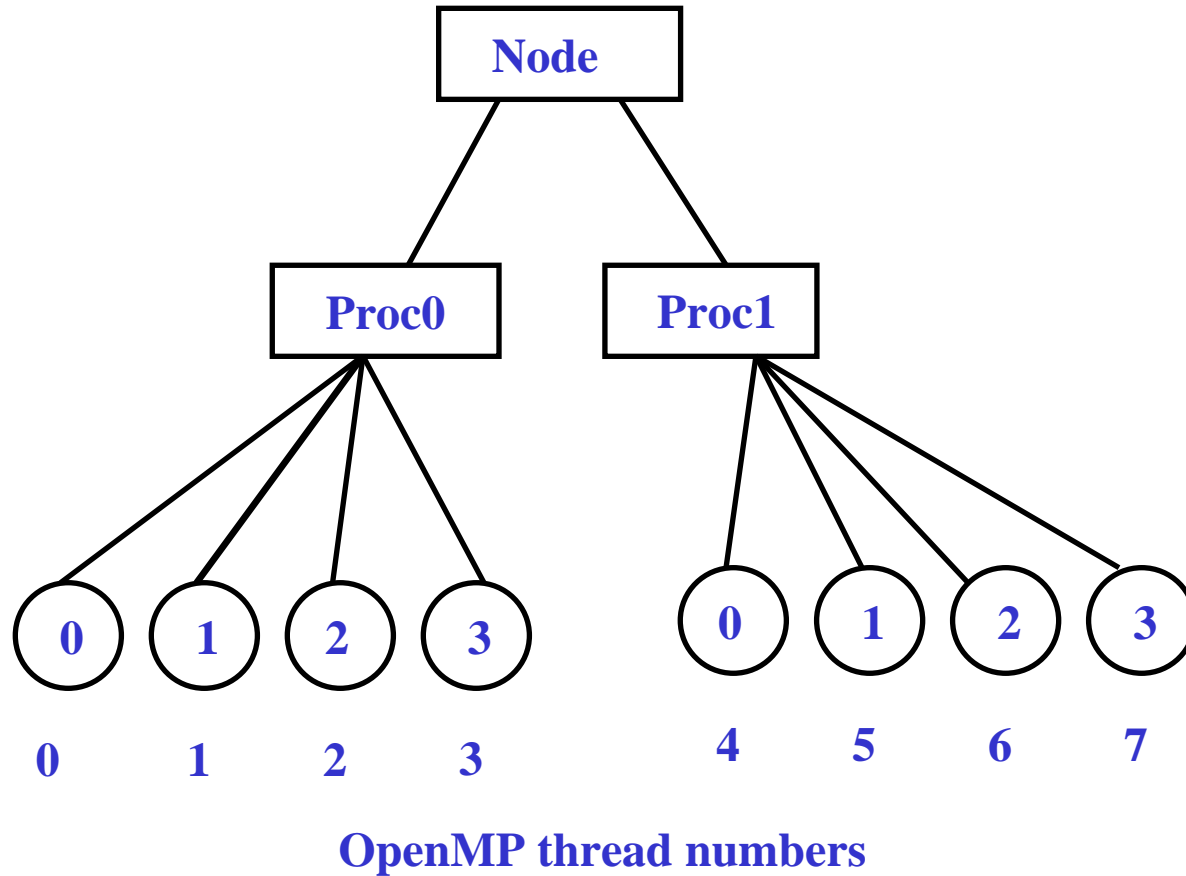
# CPU Affinity

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- Intel's OpenMP runtime library allows OpenMP threads to be bound to physical processing units
- Can have a dramatic impact on the runtime of the code
- Need to use `KMP_AFFINITY`
  - ✓ Can determine machine topology and assign OpenMP threads to processors
- On HPC-FF, `KMP_AFFINITY` overwritten by PSI daemon
- Need to export two other variables
  - ✓ `__PSI_NO_PINPROC -> 1`
  - ✓ `__PSI_NO_MEMBIND -> 1`

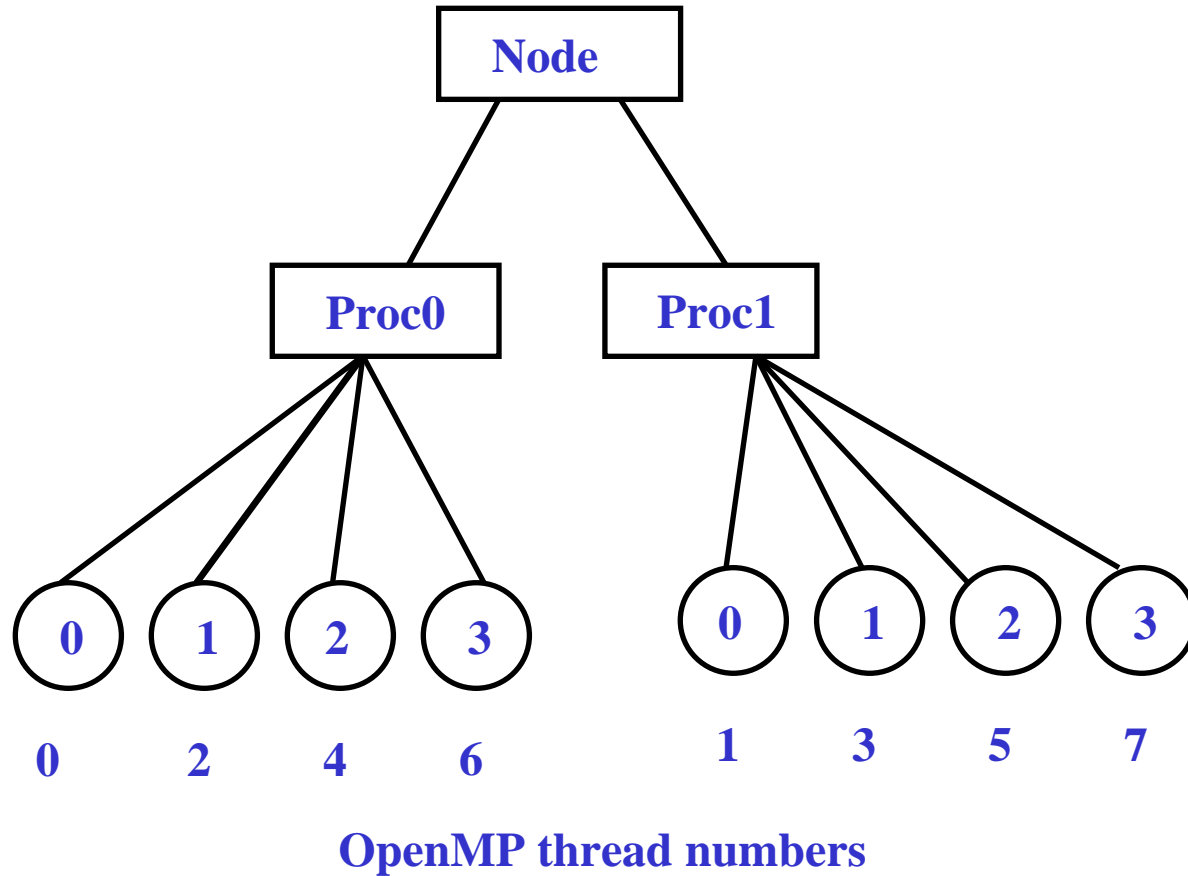


# KMP\_AFFINITY=compact





# KMP\_AFFINITY=scatter







# Additional information

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- Use module commands in the batch script
- To check correct paths of the executable, use ldd in the batch script
- Not all .a files found in the compute nodes, compile in the login node while debugging.
  - Statically linked files have a ‘.a’ extension, so an executable compiled in the login node can still run on the compute node.
  - Compiling on the compute node will not work as all the .a files are required during the compilation process.



# Summary

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- **Two versions of ParTec MPI, Intel MPI**
- **MPI consumes memory, use PSP\_ONDEMAND where ever possible**
- **Avoid MPI\_ANY\_SOURCE**
- **Use PSI\_TPP to increase available memory and bandwidth**



# Acknowledgements

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- **Juelich Support - Alexander Schnurpfeil, ParTec - Jens Hauke**
- **ORB5 – Alberto Bottino, IPP**
- **GENE – Tilman Dannert, IPP**
- **GEMR – Bruce D Scott, IPP**
- **JOREK – Guido Huysmans, CEA; Florent Sourbier (Benchmark), CEA**
- **MDCASK – Maria Jose Caturla, Univ of Alicante**
- **GYSELA - Virginie Grandgirard (Benchmark on HPC-FF), CEA**
- **FIPC module - Ian Bush, NAG**